

Advanced Microeconomics

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Exercises 2

Exercise 1 *Prove that a fully cooperative strategy profile is strongly Pareto efficient.*

Exercise 2 *Determine which of the following bimatrix games are a prisoner's dilemma.*

a. $\begin{pmatrix} 3; -1 & 3; 1 & 6; 0 \\ 1; 0 & 3; 1 & 6; 0 \\ 2; 2 & 4; 1 & 8; 2 \end{pmatrix}$.

b. $\begin{pmatrix} 1; 0 & 3; 1 & 6; 0 \\ 2; 1 & 4; 1 & 8; 1 \end{pmatrix}$.

c. $\begin{pmatrix} 6; 1 & 3; 1 & 1; 5 \\ 2; 4 & 4; 2 & 2; 3 \\ 5; 1 & 6; 1 & 5; 2 \end{pmatrix}$.

d. $\begin{pmatrix} -1; -1 & 2; 0 \\ 0; 2 & 3; 3 \end{pmatrix}$.

e. $\begin{pmatrix} 2; 2 & -1; 3 \\ 3; -1 & 0; 0 \end{pmatrix}$.

Exercise 3 *Answer the following true/false questions concerning bimatrix games.*

- a. *A bimatrix game concerns a game with two players.*
- b. *Each bimatrix game has at least one Nash equilibrium.*
- c. *Each bimatrix game has a strictly dominant strategy.*
- d. *Each bimatrix game has a fully cooperative strategy profile.*
- e. *Each bimatrix game has a weakly Pareto efficient strategy profile.*
- f. *Each fully cooperative strategy profile is weakly Pareto efficient.*
- e. *A strictly dominant strategy is fully cooperative.*
- f. *A prisoners' dilemma game has a Nash equilibrium.*
- g. *It is impossible that a weakly Pareto inefficient strategy profile is a Nash equilibrium.*
- h. *A Nash equilibrium is a strategy profile that consists of strategies of the players' that they like the most.*

Exercise 4 The following true/false questions deal with the bimatrix game

$$\begin{pmatrix} 3; 6 & 6; 5 & 7; -3 \\ -6; 2 & 5; 3 & 5; 4 \end{pmatrix}.$$

- a. The row-player has 2 strategies.
- b. There are 6 strategy profiles.
- c. The strategy profile (1, 1) is a Nash equilibrium.
- d. The row-player has a strictly dominant strategy.
- e. There is a weakly Pareto inefficient nash equilibrium.
- f. The column-player has a strictly dominant strategy.
- g. This game is a prisoners' dilemma.
- h. Playing row 1 and column 3 is a fully cooperative strategy profile
- i. This game is a zero-sum game.
- j. (1, 2) is a weakly Pareto efficient strategy profile.

Exercise 5 Consider the Hotelling game in the case $m = 2$ (so there are three vertices) and $w = 1$ (i.e. inelastic case). Determine the Nash equilibria of this game

- a. Represent this game as 3×3 -bi-matrix game with at the first row strategy 0 for player 1, at the second row strategy 1 for player 1, etc.
- b. Determine the Nash equilibria, the strongly Pareto efficient strategy profiles and the weakly Pareto efficient strategy profiles.

Exercise 6 Again consider the Hotelling Game with sites $0, 1, \dots, m$. Suppose m is even.

- a. Show that for the payoff function f_1 of player 1

$$f_1(x_1, x_2) := \begin{cases} \frac{x_1 + x_2 + 1}{2} & \text{if } x_1 < x_2, \\ \frac{m+1}{2} & \text{if } x_1 = x_2, \\ m + 1 - \frac{x_1 + x_2 + 1}{2} & \text{if } x_1 > x_2 \end{cases}$$

- b. Show that $(m/2, m/2)$ is a Nash equilibrium.

Short solutions.

Solution 1 We prove this by contradiction. So suppose \mathbf{x} is fully cooperative and \mathbf{x} would not be strongly Pareto efficient. Then there exists a pareto improvement \mathbf{y} of \mathbf{x} . In \mathbf{y} the sum of payoffs is greater than in \mathbf{x} . This is a contradiction with \mathbf{x} being fully cooperative.

Solution 2 Only the game in e is a prisoner's dilemma game.

Solution 3 aT bF cF dT eT fT gF hT iF jF.

Some explanation. Concerning f (each fully cooperative strategy profile is weakly Pareto efficient): suppose the strategy profile \mathbf{x} is fully cooperative, meaning that the total payoff is maximal. If it would not be weakly Pareto efficient, then there is a strategy profile which is better for both players and thus leads to a greater payoff than in \mathbf{x} . (In fact each fully cooperative strategy profile even is strongly Pareto efficient. In order to see this modify the above reasoning in an appropriate way.)

Concerning e: as each bimatrix game has a fully cooperative strategy profile, part f implies that each bimatrix game has a weakly Pareto efficient strategy profile.

Solution 4 aT bT cT dT eF fF gF hF iF jT.

Solution 5 a.

$$\begin{pmatrix} 3/2; 3/2 & 1; 2 & 3/2; 3/2 \\ 2; 1 & 3/2; 3/2 & 2; 1 \\ 3/2; 3/2 & 1; 2 & 3/2; 3/2 \end{pmatrix}.$$

b. There is a unique Nash equilibrium: the strategy profile (2, 3), i.e. (vertex 1, vertex 1).

c. Each strategy profile is strongly Pareto efficient and weakly Pareto efficient (and even fully cooperative).

Solution 6 a. Make a figure and count the contributions. In doing so, distinguish between $x_1 + x_2$ even and $x_1 + x_2$ odd.

b. We have to show that $f_1(x_1, m/2) \leq f_1(m/2, m/2)$ for all x_1 and that $f_2(m/2, x_2) \leq f_2(m/2, m/2)$ for all x_2 . We prove here the first statement; the second follows in the same way.

For $x_1 = m/2$, the statement is clear. For $x_1 < m/2$, we have, using part a, $f_1(x_1, m/2) = \frac{x_1 + \frac{m}{2} + 1}{2} = \frac{x_1}{2} + \frac{m}{4} + \frac{1}{2} < \frac{m}{4} + \frac{m}{4} + \frac{1}{2} = \frac{m+1}{2} = f_1(\frac{m}{2}, \frac{m}{2})$. And for $x_1 > m/2$, we have, using part a, $f_1(x_1, m/2) = m + 1 - \frac{x_1 + \frac{m}{2} + 1}{2} = m + 1 - \frac{x_1}{2} - \frac{1}{2} - \frac{m}{4} = \frac{3}{4}m + \frac{1}{2} - \frac{x_1}{2} > \frac{3}{4}m + \frac{1}{2} - \frac{m}{4} = \frac{m+1}{2} = f_1(\frac{m}{2}, \frac{m}{2})$.